

QUIESCE TERMINATION/SUSPENSION  
IN A MULTITHREADED ENVIRONMENT

Background of the Invention

Field of the Invention

5           This invention relates to a method of coordinating the quiescing (i.e., termination or suspension) of the various threads of a multithreaded process.

Description of the Related Art

10           Computer operating systems -- the software that interfaces between user applications and the hardware and performs the basic supervisory functions in a computer system -- are well known in the art. Many modern operating systems allow for the use of multiple threads within a process, or application. A multithreaded application is  
15           defined as a program using more than one thread of control to perform its work. (The terms "process" and "application" are used interchangeably in this specification to refer to one or more threads sharing a common address space.) A. S. Tanenbaum, Modern Operating Systems, (1992), incorporated  
20           herein by reference, describes several modern operating systems generally, as well as threads in particular at pp. 507-23.

25           A particular example of an operating system supporting multithreaded applications is the IBM MVS/ESA operating system with its recently introduced OpenEdition MVS extension. The OpenEdition MVS extension allows applications written to the IEEE POSIX 1003.1, 1003.2 and 1003.4a (draft) standards to run on a hardware-software platform made up of an IBM System/390 computer and the  
30           MVS/ESA operating system. (IBM, OpenEdition, MVS/ESA and

System/390 are trademarks of IBM Corporation.) Further information on the OpenEdition MVS extension may be found in the following publications, which are incorporated herein by reference:

5           Ault, "Fork Clone Address Space Implementation on MVS",  
IBM Technical Disclosure Bulletin, vol. 35, no. 6, pp.  
363-67 (Nov. 1992);

          Ault, "Interoperability Between MVS and POSIX  
10           Functions", IBM Technical Disclosure Bulletin, vol. 35,  
no. 6, pp. 383-88 (Nov. 1992);

          Ault et al., "Cross-Address Space Control Function",  
IBM Technical Disclosure Bulletin, vol. 36, no. 10, pp.  
591-95 (Oct. 1993);

15           Introducing OpenEdition MVS, IBM Publication No. GC23-  
3010-00 (1993);

MVS/ESA Support for IEEE POSIX Standards: Technical  
Presentation Guide, IBM Publication No. GG24-3867-00  
(1993).

20           As noted above, the OpenEdition MVS extension of the  
MVS/ESA operating system allows for the use of multiple  
threads within a process. In MVS terms, a thread can be  
thought of as a task. Multiple threads thus equate to the  
use of multiple MVS tasks within an MVS address space.

25           Although multithreaded applications are advantageous in  
many situations, lack of adequate task control in a  
multitasking (i.e., multithreaded) address space causes  
problems in termination, debugging and dumping. Thus, the  
POSIX standard calls for the termination of all threads  
within a process if any one of those threads terminates

abnormally. This can be accomplished in MVS by abending the job step task or by using CallRTM to abend the appropriate tasks. Many problems are encountered however, when these types of asynchronous abends are sent to the MVS tasks that  
5 were supporting OpenEdition MVS threads.

One problem that occurs is that the run-time library cannot serialize its cleanup of common process resources when the threads of the process are taken down in this abrupt, random manner. Another is that many components do  
10 not have sufficient error recovery to handle being abended between any two instructions. In some cases these deficiencies can have catastrophic results, destruction of the file system, to name one. Although the abend error recovery procedure might be improved, it would be preferable  
15 to avoid this type of abending altogether.

The desire to suspend the remaining threads of a multithreaded application in a controlled manner may arise in a debugging context. When debugging a multithreaded application, it would be desirable to allow a user debugging  
20 such an application to choose which threads run and which threads are suspended for any particular event and to be able to change the run/suspend status dynamically. This suspension process also should be of a sort that neither changes the flow of the application nor allows any thread to  
25 hold a critical system-managed resource at the time of suspension.

Another context in which the desire to suspend the remaining threads of a multithreaded application may arise is when obtaining a dump of the process with information  
30 captured from all of the threads. The desire here is similar to that in the debugging situation described above. The task requesting the dump should be able to suspend the execution of all the other tasks such that the other tasks

do not hold any critical system resources that would prevent the calling task from taking the dump. After the dump is taken, the dumping task must resume execution of the other tasks.

5           Thus, lack of adequate task control in a multitasking address space causes problems in termination, debugging and dumping. What is desired is a mechanism for terminating or suspending execution of tasks in a multithreaded environment in a predictable and nondestructive fashion.

## 10                           Summary of the Invention

          The above problems are resolved by creating a new quiesce function that when invoked sends a quiesce event to all other threads (i.e., tasks) in the address space. The thread invoking the quiesce function then waits until all  
15       the events are acted upon and the target threads are placed into the desired state.

          To accomplish this, a registration function is provided that allows a user to make known, to the operating system, the quiesce exit routine that is to be given control upon  
20       receipt of a quiesce event. If no exit is specified then the operating system determines when the event is handled.

          The method of delivery of the event is a service request block/interruption response block (SRB/IRB) combination that interrupts the target thread's execution.  
25       From here various checks can be made on the request block (RB) that was running at the time of the interrupt to see if it is safe to act upon the quiesce event. These checks include making sure that the environment is acceptable for giving control to the quiesce exit; one does not want to  
30       interrupt a system service, for instance. If control is given to the quiesce exit and it determines that the quiesce

event can be acted on, the appropriate action is taken. If the quiesce event was for termination then the exit terminates the thread. If the quiesce event was to suspend then the exit issues the appropriate suspension service.

5        If the system IRB or the user's quiesce exit determines that the quiesce event cannot be acted on, then the event is left pending and delivered again upon exit of the next system service or even sooner if the user detects that it has reached a safe point to act on the quiesce event.

10       The delivery of the quiesce event is carried out on all of the threads until the last thread has entered the desired state. The last thread doing so posts (i.e., resumes) the invoker of the quiesce function.

15       The advantage of this solution lies in allowing the decision of when the quiesce state is to be entered up to the thread being effected. The problem of unconditionally stopping a thread while holding a critical resource is avoided. Also avoided is the destructive results of asynchronously abending a thread executing in an "unstable" or "critical" section of code.

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#### Brief Description of the Drawings

Fig. 1 is a schematic block diagram of a computer system incorporating the present invention.

25       Fig. 2 shows the generic procedure for quiescing (i.e., freezing or terminating) other threads in accordance with the present invention.

Fig. 3 shows the procedure for initially registering a quiesce exit in accordance with the present invention..

Fig. 4 shows the initial phase of the procedure for freezing (i.e., suspending) other threads in accordance with the present invention.

5 Fig. 5 shows the quiesce exit interface block (QEIB) created during the phase of the freeze procedure shown in Fig. 4.

Fig. 6 shows the thread control queue element (TCQE) created during the initial registration procedure shown in Fig. 3.

10 Fig. 7 shows the final phase of the procedure for freezing other threads in accordance with the present invention.

15 Fig. 8 shows the procedure for unfreezing (i.e., resuming) other threads in accordance with the present invention.

Fig. 9 shows the procedure for terminating other threads in accordance with the present invention.

#### Description of the Preferred Embodiment

20 Fig. 1 is an overview of a computer system 100 incorporating the present invention, showing the relationships between the system layers for the implementation of the quiesce function. These layers are, starting from the top of Fig. 1, the application layer 102, the language subroutine and run-time library (RTL) layer  
25 104, the operating system or kernel layer 106, and the hardware layer 108.

At the base of the system 100 is the hardware layer 108, which consists of one or more central processing units

(CPUs), main memory and input/output (I/O) devices such as magnetic disk drives, magnetic tape drives, terminals, printers and the like. These components are conventional in the art and are hence not separately shown.

5           Kernel layer 106, the next layer above the hardware layer 108, consists of software that controls the base hardware layer, managing its CPU(s), memory and I/O devices, and contains a set of callable services (including the  
10           code access to the resources at the hardware layer. Kernel layer 106 may comprise the IBM MVS/ESA operating system with its OpenEdition MVS extension, running on a hardware layer 108 comprising an IBM System/390 computer. However, the present invention is not limited to such a hardware/software  
15           platform, and other platforms could alternatively be used.

          Language subroutine and run-time library (RTL) layer 104 is located directly above the kernel layer 106. Layer 104 consists of routines to support high-level languages (HLLs) used by many applications. Layer 104 essentially  
20           transforms the requested HLL function into the appropriate kernel service(s) to complete the request.

          Application layer 102, containing one or more user applications, is the top layer of the system 100. Applications can request system services directly from the  
25           kernel layer 106 or via a HLL through the language subroutine and run-time library layer 104. For the purposes of this specification, the application layer 102 and the language subroutine and run-time library 104 may be collectively regarded as the application.

30           Fig. 2 is a high-level flow diagram of the quiesce function of the present invention as implemented in system 100. Although the flows for specific quiesce functions

(freeze, unfreeze, or terminate) are slightly more involved, all quiesces follow this general flow.

Fig. 2 shows a user application 200, from application layer 102 of system 100 (Fig. 1), comprising a first thread 202 (thread 1) and one or more additional threads 204 (threads 2 through N). Threads 202, 204 have many of the attributes of independent processes -- e.g., they each have their own program counters and states (i.e., running, ready or blocked) -- but differ from such independent processes in that they share a common user address space 206. (Application 200 may in fact be defined as the set of threads 202, 204 sharing the same address space.) Thread 202 is not necessarily the first thread of application 200 to be created; it is distinguished from the other threads 204 merely because it invokes the quiesce services to be described.

Also shown in Fig. 2 is the quiesce service 208 of the present invention, from kernel layer 106 of system 100. Quiesce service 208 resides in a kernel address space 210 separate from the user address space and includes the following separate services, to be described below:

1. quiesce\_exit\_registration (306):

Allows the initial thread to provide the address of the quiesce exit that is to receive control when a thread requests one of the quiesce functions listed below. A quiesce exit is an application-level routine that receives control as a result of a quiesce event occurring on a thread that is associated with (created by) the initial thread that registered. This exit is responsible for acting on the quiesce event delivered.



2. `quiesce_freeze (404):`

Sends a suspension event to all other threads, giving control to the user-defined quiesce exit. For this event, the quiesce exit determines whether any critical resources are held by the thread 204. If no critical resources are held, the quiesce exit calls the `quiesce_freeze_self` service. If critical resources are held, processing of the quiesce event is delayed until the critical resources have been freed.

3. `quiesce_freeze_self (722):`

Suspends the invoking thread. This option is used from the quiesce exit on receipt of a suspension event.

4. `quiesce_event_put_back (734):`

Delays the processing of a quiesce event until a later time. This option is used from the quiesce exit on receipt of a suspension event that cannot be processed due to the current execution environment.

5. `quiesce_unfreeze (806):`

Resumes all frozen threads.

6. `quiesce_term (904):`

Sends termination events to all other threads, giving control to the user-defined quiesce exit if one was specified via the `quiesce_exit_registration` function. For this event, the

5 quiesce exit determines whether any critical  
resources are held by the thread 204. If no  
critical resources are held, the quiesce exit  
calls the pthread\_exit service to terminate the  
thread. If critical resources are held,  
processing of the quiesce event is delayed until  
the critical resources have been freed.

7. quiesce\_force:

10 Sends termination events to all other threads,  
bypassing the call to the user-defined quiesce  
exit.

15 The present invention comes into operation when an  
application event 212 is detected on thread 202. This  
application event 212 cannot be processed while other  
threads 204 in the application 200 are executing. Event 212  
could be an abend, a program check or a break point. At  
this point, application 200 may want to suspend the  
execution of the other threads 204 so that it can take a  
dump or perform diagnostic functions. Alternatively, the  
20 application 200 may want to inform the other threads 204  
that they need to terminate.

25 The present invention may be used in a debugging  
context, as described above. However, the present invention  
is not limited to such use, and the particulars of such use  
are not part of the present invention.

30 In accordance with the present invention, when thread  
202 is notified of an event 212 that requires the other  
threads 204 in the application 200 to be quiesced, it calls  
(214) the kernel quiesce service 208 to quiesce the other  
threads in the application. The quiesce service 208 sends  
(216) a quiesce notification 218 to the other threads 204

and waits (i.e., suspends) (220) until the quiesce notification has been acted upon by the other threads. Upon receiving the quiesce notification 218, threads 204 take the appropriate action based on the quiesce type. When the last application thread 204 acts on the quiesce notification, it posts (i.e., resumes) (222) thread 202, which is waiting in the kernel quiesce service 208. Upon the delivery of the post to thread 202, the quiesce service 208 returns (224) to the application 200 on thread 202.

Thread 202 can now perform (226) any critical code which could not have been done while other threads 204 were running. When this critical code is complete and the original action was to freeze the other threads 204, thread 202 calls (228) the quiesce service 208 to unfreeze the other threads.

Fig. 3 shows how application 200 registers its quiesce exit, an application-defined routine (Figs. 7-9) that handles quiesce events in a manner described further below.

To register the quiesce exit, the first thread 302 of the application 200 to be created calls (304) the quiesce\_exit\_registration service 306 of the quiesce service 208, passing the address (quiesce\_exit) of the quiesce exit routine. Quiesce\_exit\_registration service 306 stores the application's quiesce exit address in a thread control queue element (TCQE) 308 for the invoking thread 302. Referring to Fig. 6, TCQE 308 is a defined area in memory that includes a thread identifier 602 identifying the thread 302, the application quiesce exit address 604, and an event control block 606 for wait and post. Thread 302 is now registered for quiesce events.

Thread 302 can now create (310) additional threads 312 using a thread creation service 314 (pthread\_create) in the

kernel address space. Kernel thread creation service 314 may be implemented in any suitable manner known to the art; the manner in which it is implemented is not part of the present invention. For each new thread 312 that it creates  
5 (316), the thread creation service builds (318) a TCQE 308 similar to the TCQE 308 for thread 302, assigning each new thread a unique thread identifier 602 and copying the quiesce exit address 604 from the TCQE 308 of the creating thread 302 to the TCQE 308 of the newly created thread 312.  
10 Thus all threads 302, 312 in the application are registered with the same quiesce exit address 604. As shown in Fig. 3, each TCQE 308 is linked to the TCQE for the next thread 312 by any suitable means (e.g., pointers, contiguous memory locations or the like) to form a chain, or thread control  
15 queue (TCQ) 320 as it is referred to below, so that the TCQEs 308 for a given application may be scanned sequentially as described below.

Fig. 4 shows in more detail the process flow shown generally in Fig. 2, beginning with the event 212 occurring  
20 on application thread 202 that requires all other threads 204 in the application 200 to be frozen.

When the application thread 202 detects event 212, it calls (402) the quiesce\_freeze service 404 of the quiesce  
25 service 208. Quiesce\_freeze service 404 in turn invokes (406) an internal event generator 408.

For each thread 204 in the application 200 other than the invoking thread 202, the event generator 408 first creates (410) a quiesce exit interface block (QEIB) 412.  
30 This is accomplished by searching TCQ 320 and identifying all the prospective threads 204. Referring to Fig. 5, QEIB 412 is a defined area in user address space that includes locations for storing the address 502 of the target thread, the quiesce event type 504 (i.e., freeze or terminate), and

the interrupt program status word (PSW) and register contents 506 to be described below. Event generator 408 initially fills the QEIB 412 with the target thread address 502 and the quiesce event type 504.

5           For each thread 204 found in TCQ 320, event generator 408 then schedules (414) a service request block (SRB) 416 to send a freeze request interrupt to the thread, to suspend it. Each SRB 416 is a unit of work that is dispatched by the kernel to execute in the user address space 206. Each  
10       SRB 416 creates and schedules an interrupt request block (IRB) 418 to the target thread 204. IRBs 418 operate in the manner described further below.

          Once event generator 408 has scheduled the interrupts to all the appropriate threads 204, it returns (420) to the  
15       quiesce\_freeze service 404.

          Quiesce\_freeze service 404 then waits (422) on the event control block 606 (Fig. 6) located in the TCQE 308 for thread 202. This wait is satisfied by a post from the last thread 204 to respond to the quiesce event, just before the  
20       last thread enters the quiesced frozen state, as described below.

          Fig. 7 describes the interrupt mechanism in more detail. Once an IRB 418 gains control in the application 200, execution halts on the targeted thread 204 at point  
25       702. The IRB 418 then checks the application's thread status to ensure that the system environment is acceptable for handling a quiesce event. If the environment is acceptable (704), then IRB 418 saves (706) the program status word (PSW) and register contents for the thread 204  
30       in portion 506 of the corresponding QEIB 412 (Fig. 5). The PSW and register contents for the thread 204 are extracted from a corresponding task-level control block 708 which is

managed by the system dispatcher for thread 204. The PSW points to the interrupt point 702 which is to be resumed after the quiesce event is handled. IRB 418 then modifies (710) the resume PSW of the targeted thread 204 to point to the previously registered quiesce exit (712), and modifies (714) register 1 to point to the QEIB 412 of the thread 204. IRB 418 then exits (716) and execution is resumed on the target thread 204 with the modified PSW and register 1.

The quiesce exit 712 gains control and has access to the QEIB 412. The quiesce exit 712 verifies the application environment (i.e., the execution state of the thread) to make sure the thread 204 is not holding any critical resources that could deadlock the process.

If the environment is acceptable (718), the quiesce exit 712 examines the quiesce event type field 504 in QEIB 412 (Fig. 5) to determine which event type it has been called to process and, hence, whether to suspend itself or terminate. The two possible event types are an event calling for suspension of the thread 204 and an event calling for termination of the thread. If, as in this example, the event is of a type calling for suspension of the thread 204, then the quiesce exit 712 invokes (720) the quiesce\_freeze\_self service 722 of the quiesce service 208. Quiesce\_freeze\_self service 722 checks to see if the invoking thread 204 is the last thread in the application 200 to reach the quiesced state. If so (724), it posts (726) the quiesce\_freeze service 404 and hence the thread 202 (Fig. 4) that originally invoked the quiesce\_freeze service. The quiesce\_freeze\_self service 722 then suspends (728) execution of thread 204 by calling a system wait service, causing thread 204 to enter a wait.

If the quiesce exit 712 finds that the application environment is not acceptable (730), then it invokes (732)

the quiesce\_event\_put\_back service 734 to return the event back to the kernel. Quiesce\_event\_put\_back service 734 notifies the kernel that the quiesce event cannot be handled at this time and that the application 200 will request  
5 delivery of the event at a later time. This is done by marking the event control block 606 (Fig. 6) of the appropriate TCQE 308 to indicate that the quiesce event is still pending. The burden of clearing the environment of any obstacles is on the application 200. Once the  
10 application thread 204 is ready to be quiesced, it calls the quiesce\_freeze\_self service 722 or, alternatively, requests the kernel to redrive the quiesce exit 712.

As noted above, when all threads 204 have invoked the quiesce\_freeze\_self service 722, the quiesce\_freeze\_self  
15 service posts (726) the quiesce\_freeze service 404 (Fig. 4) and control is returned to its invoker (thread 202 of application 200). Referring now to Fig. 8, with all other threads 204 now frozen, the lone running thread 202 can now perform (802) the critical work desired without hindrance  
20 from the other threads 204 in the application 200. Upon completion of the critical work, the thread 202 calls (804) the quiesce\_unfreeze service 806 of the quiesce service 208. The quiesce\_unfreeze service 806 runs through the queue 320 of thread control queue elements 308 built during pthread\_  
25 create (Fig. 3) and posts (808) each thread 204 that is currently waiting inside the quiesce\_freeze\_self service 722. Upon being posted, quiesce\_freeze\_self service returns (810) control of each frozen thread 204 back to the corresponding quiesce exit 712, which resumes (812)  
30 execution at the point of interruption 702; this is done by using the PSW and register contents 506 saved in the QEIB 412. The quiesce\_unfreeze service 806 then returns (814) control to thread 202, which resumes (816) application processing.

Fig. 9 shows the flow of a quiesce termination request. A quiesce termination request results in the termination of all threads 204 in the application 200, with the exception of the calling thread 202. The terminating threads 204 are given an opportunity to complete application-critical code and/or clean up thread-related resources before being terminated.

As shown in Fig. 9, upon being notified of an event 212 that requires the other threads 204 in the application 200 to be terminated, thread 202 calls (902) the quiesce\_term service 904 of the quiesce service 208 to request termination of the other threads in the application. The quiesce\_term service 904 scans the TCQ chain 320 and generates (906) a quiesce\_term event 908 for each TCQE 308 on the chain, excluding the one for the calling thread 202. Although not shown in Fig. 9, this quiesce\_term event 908 is preferably delivered to threads 204 using service request blocks (SRBs) and interrupt request blocks (IRBs) similar to those shown in Fig. 4. The quiesce\_term service 904 then waits (910) while the other threads 204 are terminated.

When threads 204 receive the event 908 generated by the quiesce\_term service 904, the normal flow of the application is interrupted at point 702, as before, and the quiesce exit 712 is given control. The quiesce exit 712 checks the application environment, as in the suspension case (Fig. 7) described above. If the application environment is acceptable, and the generated quiesce event (as indicated by the quiesce event type field 504 in QEIB 412) is for termination (912), then the quiesce exit 712 invokes (914) the thread cleanup routine (pthread\_exit) 922.

The pthread\_exit routine 922 releases system resources associated with the terminating thread 204. If the terminating thread 204 is the last thread in the application



that had a quiesce\_term event generated to it (916), then  
the pthread\_exit routine 922 posts (918) thread 202, which  
is waiting in the quiesce\_term service 904. After thread  
202 is posted (918) out of its wait in quiesce\_term service  
5 904, it returns (920) to the application 200 to take the  
appropriate action based on the event 212 received.

The quiesce\_force service (not shown), operates in a  
manner similar to that of quiesce\_term service 904, sending  
termination events to all other threads 204. However, the  
10 quiesce\_force service bypasses the call to the user-defined  
quiesce exit, and the pthread\_exit routine 922 is called  
from the interrupt request block (IRB) directly.

What is claimed is: